Abstract—We describe a noncooperative interference alignment (IA) technique which allows an opportunistic multiple input multiple output (MIMO) link (secondary) to harmlessly coexist with another MIMO link (primary) in the same frequency band. Assuming perfect channel knowledge at the primary receiver and transmitter, capacity is achieved by transmitting along the spatial directions (SD) associated with the singular values of its channel matrix using a water-filling power allocation (PA) scheme. Often, power limitations lead the primary transmitter to leave some of its SD unused. Here, it is shown that the opportunistic link can transmit its own data if it is possible to align the interference produced on the primary link with such unused SDs. We provide both a processing scheme to perform IA and a PA scheme which maximizes the transmission rate of the opportunistic link. The asymptotes of the achievable transmission rates of the opportunistic link are obtained in the regime of large numbers of antennas. Using this result, it is demonstrated that depending on the signal-to-noise ratio and the number of transmit and receive antennas of the primary and opportunistic links, both systems can achieve transmission rates of the same order.

Index Terms—Channel eigenmodes, cognitive radio, MIMO interference channel, opportunistic interference alignment, water-filling.

I. INTRODUCTION

The concept of cognitive radio is well known by now. The main idea is to let a class of radio devices, called secondary systems, opportunistically access certain portions of spectrum previously used by other radio devices, called primary systems, at a given time or geographical area [2]. These pieces of unused spectrum, known as white-spaces, appear mainly when either transmissions in the primary network are sporadic, i.e., there are periods over which no transmission takes place, or there is no network infrastructure for the primary system in a given area, for instance, when there is no primary network coverage in a certain region. In the case of dense networks, a white-space might be a rare and short-lasting event. In fact, the idea of cognitive radio as presented in [2] (i.e., spectrum pooling), depends on the existence of such white-spaces [3]. In the absence of those spectrum holes, secondary systems are unable to transmit without producing additional interference on the primary systems. One solution to this situation has been provided recently under the name of interference alignment (IA). Basically, IA refers to the construction of signals such that the resulting interference signal lies in a subspace orthogonal to the one spanned by the signal of interest at each receiver. The IA concept was independently introduced by several authors [4]–[7]. Recently, IA has become an important tool to study the interference channel, namely its degrees of freedom [6], [8], [9]. The feasibility and implementation issues of IA regarding mainly the required channel state information (CSI) has been also extensively studied [10]–[13].

In this paper we study an IA scheme named opportunistic IA (OIA) [1]. The idea behind OIA can be briefly described as follows. The primary link is modeled by a single-user MIMO channel since it must operate free of any additional interference produced by secondary systems. Then, assuming perfect CSI at both transmitter and receiver ends, capacity is achieved by implementing a water-filling power allocation (PA) scheme [14] over the spatial directions associated with the singular values of its channel transfer matrix. Interestingly, even if the primary transmitters maximize their transmission rates, power limitations generally lead them to leave some of their spatial directions (SD) unused. The unused SD can therefore be reused by another system operating in the same frequency band. Indeed, an opportunistic transmitter can send its own data to its respective receiver by processing its signal in such a way that the interference produced on the primary link impairs only the unused SDs. Hence, these spatial resources can be very useful for a secondary system when the available spectral resources are fully exploited over a certain period in a geographical area. The idea of OIA, as described above, was first introduced in [1] considering a very restrictive scenario, e.g., both primary and secondary devices have the same number of antennas and same power budget. In this paper, we consider a more general framework where devices have different number of antennas, different
power budgets and no conditions are imposed over the channel
transfer matrices (In [1], full rank condition was imposed over
certain matrices).

The rest of this paper is structured as follows. First, the system
model, which consists of an interference channel with MIMO
links, is introduced in Section II. Then, our aim in Section III is
twofold. First, an analysis of the feasibility of the OIA scheme
is provided. For this purpose, the existence of transmit oppor-
tunities (SD left unused by the primary system) is studied. The
average number of transmit opportunities is expressed as a func-
tion of the number of antennas at both the primary and sec-
ondary terminals. Second, the proposed interference alignment
technique and power allocation (PA) policy at the secondary
transmitter are described. In Section IV-B, tools from random
matrix theory for large systems are used to analyze the achieve-
able transmission rate of the opportunistic transmitter when no
optimization is performed over its input covariance matrix. We
illustrate our theoretical results by simulations in Section V.
Therein, it is shown that our approach allows the secondary link
to achieve transmission rates of the same order as those of the
primary link. Finally, in Section VI we state our conclusions and
provide possible extensions of this work.

II. SYSTEM MODEL

Notations: In the sequel, matrices and vectors are respec-
tively denoted by boldface upper case symbols and boldface
lower case symbols. An \( N \times K \) matrix with ones on its main
diagonal and zeros on its off-diagonal entries is denoted by \( \mathbf{1}_{N \times K} \),
while the identity matrix of size \( N \) is simply denoted by \( \mathbf{1}_N \).
An \( N \times K \) matrix with zeros in all its entries (null matrix) is
denoted by \( \mathbf{0}_{N \times K} \). Matrices \( \mathbf{X}^T \) and \( \mathbf{X}^H \) are the transpose and
Hermitian transpose of matrix \( \mathbf{X} \), respectively. The determin-
ant of matrix \( \mathbf{X} \) is denoted by \( |\mathbf{X}| \). The expectation operator is
denoted by \( \mathbb{E}[\cdot] \). The indicator function associated with a given set \( \mathcal{A} \) is
denoted by \( 1_{\mathcal{A}}(\cdot) \), and defined by \( 1_{\mathcal{A}}(x) = 1 \) (respectively, \( 0 \)) if \( x \in \mathcal{A} \) (respectively, \( x \notin \mathcal{A} \)). The Heaviside step function
and the Dirac delta function are respectively denoted by \( \mu(\cdot) \)
and \( \delta(\cdot) \). The symbols \( \mathbb{N}, \mathbb{R}, \) and \( \mathbb{C} \) denote the sets of nonnegative
integers, real numbers, and complex numbers, respectively.
The subsets \([0, +\infty[ \) and \( ]-\infty, 0] \) are denoted by \( \mathbb{R}^+ \) and \( \mathbb{R}^- \),
respectively. The operator \( (x)^+ \) for \( x \in \mathbb{R} \) is equivalent to the operation \( \max(0, x) \). Let \( \mathbf{A} \) be an \( n \times n \) square matrix
with real eigenvalues \( \lambda_{A,1}, \ldots, \lambda_{A,n} \). We define the empirical
eigenvalue distribution of \( \mathbf{A} \) by \( F^{(n)}(\cdot) \triangleq \frac{1}{n} \sum_{i=1}^{n} \mu(\lambda - \lambda_{A,i}) \),
and, when it exists, we denote \( f^{(n)}_{\mathbf{A}}(\lambda) \) the associated eigenvalue
probability density function, where \( F^{(n)}(\cdot) \) and \( f^{(n)}(\cdot) \) are respec-
tively the associated limiting eigenvalue distribution and prob-
ability density function when \( n \to +\infty \).

We consider two unidirectional links simultaneously oper-
ating in the same frequency band and producing mutual inter-
ference as shown in Fig. 1. The first transmitter-receiver pair
\((\text{Tx}_1, \text{Rx}_1)\) is the primary link. The pair \((\text{Tx}_2, \text{Rx}_2)\) is an
opportunistic link subject to the strict constraint that the primary
link must transmit at a rate equivalent to its single-user capacity.
Denote by \( N_i \) and \( M_i \), with \( i = 1 \) (respectively, \( i = 2 \)), the
number of antennas at the primary (respectively, secondary) re-
ceiver and transmitter, respectively. Each transmitter sends in-
dependent messages only to its respective receiver and no co-
operation between them is allowed, i.e., there is no message exchange
between transmitters. This scenario is known as the
MIMO interference channel (IC) [15], [16] with private
messages. A private message is a message from a given source to a
given destination: only one destination node is able to decode it.
Indeed, we do not consider the case of common messages which
would be generated by a given source in order to be decoded by
several destination nodes.

In this paper, we assume the channel transfer matrices be-
tween different nodes to be fixed over the whole duration of
the transmission. The channel transfer matrix from transmitter
\( j = \{1, 2\} \) to receiver \( i = \{1, 2\} \) is an \( N_i \times M_j \) matrix
denoted by \( \mathbf{H}_{ij} \) which corresponds to the realization of a random
matrix with independent and identically distributed (i.i.d.) com-
plex Gaussian circularly symmetric entries with zero mean and
variance \( 1/M_j \), which implies

\[
\mathbb{V}(i, j) \in \{1, 2\}^2, \quad \text{Trace}\left( \mathbb{E}\left[ \mathbf{H}_{ij} \mathbf{H}_{ij}^H \right] \right) = N_i.
\] (1)

The \( L_i \) symbols transmitted by transmitter \( i \) is able to simultane-
umously transmit, denoted by \( s_{i,1}, \ldots, s_{i,L_i} \), are represented by the
vector \( \mathbf{s}_i = (s_{i,1}, \ldots, s_{i,L_i})^T \). We assume that \( \forall i \in \{1, 2\} \)
symbols \( s_{i,1}, \ldots, s_{i,L_i} \) are i.i.d. zero-mean circularly-sym-
metric complex Gaussian variables. In our model, transmitter
\( i \) processes its symbols using a matrix \( \mathbf{V}_i \) to construct its
transmitted signal \( \mathbf{V}_i \mathbf{s}_i \). Therefore, the matrix \( \mathbf{V}_i \) is called
pre-processing matrix. Following a matrix notation, the primary
and secondary received signals, represented by the \( N_i \times 1 \)
column-vectors \( \mathbf{r}_i \), with \( i \in \{1, 2\} \), can be written as

\[
\begin{pmatrix}
\mathbf{r}_1 \\
\mathbf{r}_2
\end{pmatrix} = \begin{pmatrix}
\mathbf{H}_{11} & \mathbf{H}_{12} \\
\mathbf{H}_{21} & \mathbf{H}_{22}
\end{pmatrix}\begin{pmatrix}
\mathbf{V}_1 \mathbf{s}_1 \\
\mathbf{V}_2 \mathbf{s}_2
\end{pmatrix} + \begin{pmatrix}
\mathbf{n}_1 \\
\mathbf{n}_2
\end{pmatrix}.
\] (2)

where \( \mathbf{n}_i \) is an \( N_i \)-dimensional vector representing noise ef-
fects at receiver \( i \) with entries modeled by an additive white
Gaussian noise (AWGN) process with zero mean and variance
\( \sigma_i^2 \), i.e., \( \mathbf{n}_i \in \{1, 2\}, \mathbb{E}[\mathbf{n}_i \mathbf{n}_i^H] = \sigma_i^2 \mathbf{I}_{N_i} \). At transmitter \( i \in \{1, 2\} \), the \( L_i \times L_i \) power allocation matrix \( \mathbf{P}_i \) is defined by
the input covariance matrix \( \mathbf{P}_i = \mathbb{E}[\mathbf{s}_i \mathbf{s}_i^H] \). Note that sym-
bols \( s_{i,1}, \ldots, s_{i,L_i} \), \( \forall i \in \{1, 2\} \) are mutually independent and
zero-mean, thus, the PA matrices can be written as diagonal
matrices, i.e., \( \mathbf{P}_i = \text{diag}(p_{i,1}, \ldots, p_{i,L_i}) \). Choosing \( \mathbf{P}_i \) therefore

![Fig. 1. Two-user MIMO interference channel.](image-url)
means selecting a given PA policy. The power constraints on the transmitted signals $V_iS_i$ can be written as

$$\forall i \in \{1, 2\}, \quad \text{Trace}(V_iP_iV_i^H) \leq M_ip_{i,\text{max}}. \quad (3)$$

Here, we have assumed that the i.i.d. entries of matrices $H_{ij}$, for all $(i, j) \in \{1, 2\}^2$, are Gaussian random variables with zero mean and variance $1/M_j$. This assumption together with the power constraints in (3) is equivalent to considering a system where the entries of matrices $H_{ij}$ for all $(i, j) \in \{1, 2\}^2$ are Gaussian random variables with zero mean and unit variance, and the transmitted signal $V_iS_i$ are constrained by a finite transmit power $p_{i,\text{max}}$. Nonetheless, the first assumption allows us to increase the dimension of the system (number of antennas) while maintaining the same average received signal-to-noise ratio (SNR) level $p_{i,\text{max}}/\sigma_i^2, \forall i \in \{1, 2\}$. Moreover, most of the tools from random matrix theory used in the asymptotic analysis of the achievable data rate of the opportunistic link in Section IV-B, require the variance of the entries of channel matrices to be normalized by their size. That is the reason why the normalized model, i.e., channel transfer matrices and power constraints respectively satisfying (1) and (3), was adopted.

At receiver $i \in \{1, 2\}$, the signal $r_i$ is processed using an $N_i \times N_i$ matrix $D_i$ to form the $N_i$-dimensional vector $y_i = D_ir_i$. All along this paper, we refer to $D_i$ as the post-processing matrix at receiver $i$. Regarding channel knowledge assumptions at the different nodes, we assume that the primary terminals (transmitter and receiver) have perfect knowledge of the matrix $H_{i1}$ while the secondary terminals have perfect knowledge of all channel transfer matrices $H_{ij}, \forall (i, j) \in \{1, 2\}^2$. One might ask whether this setup is highly demanding in terms of information assumptions. In fact, there are several technical arguments making this setup relatively realistic: a) in some contexts channel reciprocity can be exploited to acquire CSI at the transmitters; b) feedback channels are often available in wireless communications [11]; and c) learning mechanisms [12] can be exploited to iteratively learn the required CSI. In any case, the perfect information assumptions provide us with an upper bound on the achievable transmission rate for the secondary link.

III. INTERFERENCE ALIGNMENT STRATEGY

In this section, we describe how both links introduced in Section II can simultaneously operate under the constraint that no additional interference is generated by the opportunistic transmitter on the primary receiver. First, we revisit the transmitting scheme implemented by the primary system [14], then we present the concept of transmit opportunity, and finally we introduce the proposed opportunistic IA technique.

A. Primary Link Performance

According to our initial assumptions (Section II) the primary link must operate at its highest transmission rate in the absence of interference. Hence, following the results in [14] and [17] and using our own notation, the optimal pre-processing and post-processing schemes for the primary link are given by the following theorem.

Theorem 1: Let $H_{11} = U_{H_{11}}A_{H_{11}}V_{H_{11}}^H$ be a singular value decomposition (SVD) of the $N_1 \times M_1$ channel transfer matrix $H_{11}$, with $U_{H_{11}}$ and $V_{H_{11}}$, two unitary matrices with dimension $N_1 \times N_1$ and $M_1 \times M_1$, respectively, and $A_{H_{11}}$ an $N_1 \times M_1$ matrix with main diagonal $(\lambda_{H_{11,1}}, \ldots, \lambda_{H_{11,n}(N_1,M_1)})$ and zeros on its off-diagonal. The primary link achieves capacity by choosing $V_1 = V_{H_{11}}, D_1 = U_{H_{11}}^{-1}, P_1 = \text{diag}(p_{1,1}, \ldots, p_{1,M_1})$, where

$$\forall n \in \{1, \ldots, M_1\}, \quad p_{1,n} = \left(1 - \frac{\sigma_i^2}{\lambda_{H_{11},n}^2}\right)^+ \quad (4)$$

with $A_{H_{11},1} = \lambda_{H_{11},1}, A_{H_{11}} = \text{diag}(\lambda_{H_{11},1}, \ldots, \lambda_{H_{11},M_1}), \lambda_{H_{11},n}$ and the constant $\beta$ (water-level) is set to saturate the power constraint (3).

Let $N \overset{\Delta}{=} \min(N_1, M_1)$. When implementing its capacity-achieving transmission scheme, the primary transmitter allocates its transmit power over an equivalent channel $D_iH_{11}V_i = A_{H_{11}}$ which consists of at most $\text{rank}(H_{11}^H) \leq N$ parallel sub-channels with nonzero channel gains $\lambda_{H_{11},n}^{-1}$, respectively. These nonzero channel gains to which we refer as transmit dimensions, correspond to the nonzero eigenvalues of matrix $H_{11}^H$. The transmit dimension $n \in \{1, \ldots, M_1\}$ is said to be used by the primary transmitter if $p_{1,n} > 0$. Interestingly, (4) shows that some of the transmit dimensions can be left unused. Let $m_1 \in \{1, \ldots, M_1\}$ denote the number of transmit dimensions used by the primary user:

$$m_1 \overset{\Delta}{=} \sum_{n=1}^{M_1} 1_{p_{1,n} \geq 0}(p_{1,n})$$

$$= \sum_{n=1}^{M_1} 1_{\lambda_{H_{11},n}^2 > 0} \left(\lambda_{H_{11},n}^2\right)^{+} \quad (5)$$

As $p_{1,\text{max}} > 0$, the primary link transmits at least over dimension $n^* = \arg\max_{m \in \{1, \ldots, \min(N_1, M_1)\}} \left\{\lambda_{H_{11},m}^2\right\}$ regardless of its SNR, and moreover, there exist at most $N$ transmit dimensions, thus

$$1 \leq m_1 \leq \text{rank}(H_{11}^H) \leq N. \quad (6)$$

In the following subsection, we show how those unused dimensions of the primary system can be seen by the secondary system as opportunities to transmit.

B. Transmit Opportunities

Once the PA matrix is set up following Theorem 1, the primary equivalent channel $D_iH_{11}V_iP_1^{1/2} = A_{H_{11}}P_1^{1/2}$ is an $N_1 \times M_1$ diagonal matrix whose main diagonal contains $m_1$ nonzero entries and $N - m_1$ zero entries. This equivalent channel transforms the set of $m_1$ used and $M_1 - m_1$ unused transmit dimensions into a set of $m_1$ receive dimensions containing a noisy version of the primary signal, and a set of $N_1 - m_1$ unused receive dimensions containing no primary signal. The $m_1$ used dimensions are called primary reserved dimensions, while the remaining $N_1 - m_1$ dimensions are named secondary transmit opportunities (TO). The IA strategy, described in Section III-C allows the secondary user to exploit these $N_1 - m_1$ receive dimensions left unused by the primary
link, while avoiding to interfere with the $m_1$ receive dimensions used by the primary link.

**Definition 2 (Transmit Opportunities):** Let $\lambda_{H_{11}} \geq \lambda_{H_{12}} \geq \cdots$ be the eigenvalues of matrix $H_{11}$, and $\beta$ be the water-level in (Theorem 1). Let $m_1$, as defined in (5), be the number of primary reserved dimensions. Then the number of transmit opportunities $S$ available to the opportunistic terminal is given by

$$S \triangleq N_1 - m_1 = N_1 - \sum_{i=1}^{M_1} \max \left[ \left( \lambda_{H_{11}} \right)_{i,i}, \left( \lambda_{H_{12}} \right)_{i,i} \right].$$

Note that in this definition it is implicitly assumed that the number of TOs is constant over a duration equal to the channel coherence time. Combining (6) and (7) yields the bounds on the number of transmit opportunities

$$N_1 - N \leq S \leq N_1 - 1.$$  

A natural question arises as to whether the number of TOs is sufficiently high for the secondary link to achieve a significant transmission rate. In order to provide an element of response to this question, a method to find an approximation of the number of TOs per primary transmit antenna, $S_{\infty}$, is proposed in Section IV-A. In any case, as we shall see in the next subsection, to take advantage of the TOs described here, a specific signal processing scheme is required in the secondary link.

**C. Pre-Processing Matrix**

In this subsection, we define the interference alignment condition to be met by the secondary transmitter and determine a pre-processing matrix satisfying this condition.

**Definition 3 (IA Condition):** Let $H_{11} = U_{H_{11}} P_{H_{11}} V_{H_{11}}^H$ be an SVD of $H_{11}$ and

$$R = \sigma_1^2 I_{N_1} + U_{H_{11}}^H V_{H_{12}} V_{H_{22}} U_{H_{12}}^H \sigma_1$$

be the covariance matrix of the co-channel interference (CCI) plus noise signal in the primary link. The opportunistic link is said to satisfy the IA condition if its opportunistic transmission is such that the primary link achieves the transmission rate of the equivalent single-user system, which translates mathematically as

$$\log_2 \left[ I_{N_1} + \frac{1}{\sigma_1^2} A_{H_{11}} P_{H_{11}} A_{H_{11}}^H \right] = \log_2 \left[ I_{N_1} + R^{-1} A_{H_{11}} P_{H_{11}} A_{H_{11}}^H \right].$$

Our objective is first to find a pre-processing matrix $V_2$ that satisfies the IA condition and then, to tune the PA matrix $P_2$ and post-processing matrix $D_2$ in order to maximize the transmission rate for the secondary link.

**Lemma 1 (Pre-Processing Matrix $V_2$):** Let $H_{11} = U_{H_{11}} A_{H_{11}} V_{H_{11}}^H$ be an ordered SVD of $H_{11}$, with $U_{H_{11}}$ and $V_{H_{11}}$ two unitary matrices of size $N_1 \times N_1$ and $M_1 \times M_1$, respectively, and $A_{H_{11}}$ an $N_1 \times M_1$ matrix with diagonal

$$A_{H_{11}} = \left[ \lambda_{H_{11,1}}, \ldots, \lambda_{H_{11,\min(N_1,M_1)}} \right]$$

and zeros on its off-diagonal, such that $\lambda_{H_{11,1}} \geq \lambda_{H_{11,2}} \geq \cdots \geq \lambda_{H_{11,\min(N_1,M_1)}}$. Let also the $N_1 \times M_2$ matrix $H \triangleq U_{H_{11}}^H V_{H_{12}}$ have a block structure,

$$H = \begin{bmatrix} M_2 \\ N_1 - m_1 \end{bmatrix} \begin{bmatrix} \tilde{H}_1 \\ \tilde{H}_2 \end{bmatrix}. \quad (11)$$

The IA condition (Def. 3) is satisfied independently of the PA matrix $P_2$, when the pre-processing matrix $V_2$ satisfies the condition:

$$\tilde{H}_1 V_2 = 0_{m_1 \times L_2} \quad (12)$$

where $L_2$ is the dimension of the null space of matrix $\tilde{H}_1$.

**Proof:** See Appendix A.

Another solution to the IA condition was given in [1], namely $V_2 = H_{12}^* U_{H_1} \tilde{P}_1$ for a given diagonal matrix $\tilde{P}_1 = \text{diag}(\tilde{\rho}_1, \ldots, \tilde{\rho}_{M_1})$, with $\tilde{\rho}_n = (\sigma_2^2 + \lambda_{H_{12}}^{H_{11,n}} \beta)^+$. Where $\beta$ is the water-level of the primary system (Theorem 1) and $n \in \{1, \ldots, M_1\}$. However, such a solution is more restrictive than (12) since it requires $H_{12}$ to be invertible and does not hold for the case when $N_1 \neq M_j, \forall (i,j) \in \{1,2\}^2$.

Plugging $V_2$ from (12) into (9) shows that to guarantee the IA condition (3), the opportunistic transmitter has to avoid interfering with the $m_1$ dimensions used by the primary transmitter. That is the reason why we refer to our technique as OIA: interference from the secondary user is made orthogonal to the $m_1$ receive dimensions used by the primary link. This is achieved by aligning the interference from the secondary user with the $N_1 - m_1$ nonuse receive dimensions of the primary link.

From Lemma 1, it appears that the $L_2$ columns of matrix $V_2$ have to belong to the null space $\text{Ker}(\tilde{H}_1)$ of $\tilde{H}_1$ and therefore to the space spanned by the $\text{dim Ker}(\tilde{H}_1) = M_2 - \text{rank}(\tilde{H}_1)$ last columns of matrix $V_{\tilde{H}_1}$, where $\tilde{H}_1 = U_{\tilde{H}_1} A_{\tilde{H}_1} V_{\tilde{H}_1}^H$ is an SVD of $\tilde{H}_1$, with $U_{\tilde{H}_1}$ and $V_{\tilde{H}_1}$ two unitary matrices of respective sizes $m_1 \times m_1$ and $M_2 \times M_2$, and $A_{\tilde{H}_1}$ an $m_1 \times M_2$ matrix containing the vector $(\lambda_{\tilde{H}_{1,1}}^{H_{1,1}}, \ldots, \lambda_{\tilde{H}_{1,\min(m_1,M_2)}}^{H_{1,1}})$ on its main diagonal and zeros on its off-diagonal, such that $\lambda_{\tilde{H}_{1,1}}^{H_{1,1}} \geq \cdots \geq \lambda_{\tilde{H}_{1,\min(m_1,M_2)}}^{H_{1,1}}$, i.e.,

$$V_2 \in \text{Span} \left( \begin{bmatrix} v_{\tilde{H}_1}^{(\text{rank}(\tilde{H}_1)+1)} \\ \vdots \\ v_{\tilde{H}_1}^{(M_2)} \end{bmatrix} \right). \quad (13)$$

Here, for all $i \in \{1, \ldots, M_2\}$, the column vector $v_{\tilde{H}_1}^{(i)}$ represents the $i$-th column of matrix $V_{\tilde{H}_1}$ from the left to the right.

In the following, we assume that the $L_2$ columns of the matrix $V_2$ form an orthonormal basis of the corresponding subspace (13), and thus, $V_{\tilde{H}_1}^H V_2 = L_2$. Moreover, recalling that $\tilde{H}_1$ is of size $m_1 \times M_2$, we would like to point out the following.

- When $m_1 < M_2$, $\text{rank}(\tilde{H}_1) \leq m_1$ and $\text{dim Ker}(\tilde{H}_1) \geq M_2 - m_1$ with equality if and only if $\tilde{H}_1$ is full row-rank. This means that there are always at least $M_2 - m_1 > 0$ nonnull orthogonal vectors in $\text{Ker}(\tilde{H}_1)$, and thus, $L_2 = \text{dim Ker}(\tilde{H}_1)$. Consequently, $V_2$ can always be chosen to be different from the null matrix $0_{M_2 \times L_2}$.
When, \( M_2 \leq m_1 \), \( \text{rank}(\mathbf{H}_1) \leq M_2 \) and \( \text{dim Ker}(\mathbf{H}_1) \geq 0 \), with equality if and only if \( \mathbf{H}_1 \) is full column-rank. This means that there are nonzero vectors in \( \text{Ker}(\mathbf{H}_1) \) if and only if \( \mathbf{H}_1 \) is not full column-rank. Consequently, \( \mathbf{V}_2 \) is a nonzero matrix if and only if \( \mathbf{H}_1 \) is not full column-rank, and again \( I_2 = \text{dim Ker}(\mathbf{H}_1) \).

Therefore, the rank of \( \mathbf{V}_2 \) is given by \( I_2 = \text{dim Ker}(\mathbf{H}_1) \leq M_2 \), and it represents the number of transmit dimensions on which the secondary transmitter can allocate power without affecting the performance of the primary user. The following lower bound on \( I_2 \) holds:

\[
I_2 = \text{dim Ker}(\mathbf{H}_1) = M_2 - \text{rank}(\mathbf{H}_1) \\
\geq M_2 - \min(M_2, m_1) \\
= \max(0, M_2 - m_1).
\]

(14)

Note that by processing \( \mathbf{s}_2 \) with \( \mathbf{V}_2 \) the resulting signal \( \mathbf{V}_2 \mathbf{s}_2 \) becomes orthogonal to the space spanned by a subset of \( m_1 \) rows of the cross-interference channel matrix \( \mathbf{H} = \mathbf{H}_1^H \mathbf{H}_2 \). This is the main difference between the proposed OIA technique and the classical zero-forcing beamforming (ZFBF) [18], for which the transmit signal must be orthogonal to the whole row space of matrix \( \mathbf{H} \). In the ZFBF case, the number of transmit dimensions, on which the secondary transmitter can allocate power without affecting the performance of the primary user, is given by \( I_{2,BF} = \text{dim Ker}(\mathbf{H}) = M_2 - \text{rank}(\mathbf{H}) \). Since \( \text{rank}(\mathbf{H}_1) \leq \text{rank}(\mathbf{H}) \), we have \( I_{2,BF} \leq I_2 \). This inequality, along with the observation that \( \text{Ker}(\mathbf{H}) \subset \text{Ker}(\mathbf{H}_1) \), shows that any opportunity to use a secondary transmit dimension provided by ZFBF is also provided by OIA, thus OIA outperforms ZFBF. In the next subsection we tackle the problem of optimizing the post-processing matrix \( \mathbf{D}_2 \) to maximize the achievable transmission rate for the opportunistic transmitter.

\[
\text{D. Post-Processing Matrix}
\]

Once the pre-processing matrix \( \mathbf{V}_2 \) has been adapted to perform IA according to (13), no harmful interference impairs the primary link. However, the secondary receiver undergoes the co-channel interference (CCI) from the primary transmitter. Then, the joint effect of the CCI and noise signals can be seen as a colored Gaussian noise with covariance matrix\( Q = \mathbf{H}_2 \mathbf{V}_{11} \mathbf{V}_1^{H} \mathbf{H}_{11}^{H} \mathbf{H}_2^{H} + \sigma_n^2 \mathbf{I}_{N_2} \).

(15)

We recall that the opportunistic receiver has full CSI of all channel matrices, i.e., \( \mathbf{H}_{11}, \mathbf{V}_{11} \) \( \subseteq \{1,2\} \). Given an input covariance matrix \( \mathbf{P}_2 \), the mutual information between the input signal \( \mathbf{s}_2 \) and the output \( \mathbf{y}_2 = \mathbf{D}_2 \mathbf{p}_2 \) is

\[
\begin{align*}
R_2(\mathbf{P}_2, \sigma_n^2) &= \log_2 |\mathbf{I}_{N_2} + D_2 \mathbf{H}_2 \mathbf{V}_{22} \mathbf{V}_2^{H} \mathbf{H}_2^{H} \mathbf{D}_2^{H} (\mathbf{D}_2 \mathbf{Q} \mathbf{D}_2^{H})^{-1}| \\
&\leq \log_2 |\mathbf{I}_{N_2} + Q^{-\frac{1}{2}} \mathbf{H}_2 \mathbf{V}_{22} \mathbf{V}_2^{H} \mathbf{H}_2^{H} Q^{-\frac{1}{2}}|.
\end{align*}
\]

(16)

where equality is achieved by a whitening post-processing filter \( \mathbf{D}_2 = Q^{-(1/2)} \) [19], i.e., the mutual information between the transmitted signal \( \mathbf{s}_2 \) and \( \mathbf{p}_2 \), is the same as that between \( \mathbf{s}_2 \) and \( \mathbf{y}_2 = \mathbf{D}_2 \mathbf{p}_2 \). Note also that expression (16) is maximized by a zero-mean circularly-symmetric complex Gaussian input \( \mathbf{s}_2 \) [14].

\[
\text{E. Power Allocation Matrix Optimization}
\]

In this subsection, we are interested in finding the input covariance matrix \( \mathbf{P}_2 \) which maximizes the achievable transmission rate for the opportunistic link, \( R_2(\mathbf{P}_2, \sigma_n^2) \) assuming that both matrices \( \mathbf{V}_2 \) and \( \mathbf{D}_2 \) have been set up as discussed in Sections III-C and III-D, respectively. More specifically, the problem of interest in this subsection is

\[
\begin{align*}
\max_{\mathbf{P}_2} & \quad \log_2 |\mathbf{I}_{N_2} + \mathbf{Q}^{-\frac{1}{2}} \mathbf{H}_2 \mathbf{V}_{22} \mathbf{V}_2^{H} \mathbf{H}_2^{H} \mathbf{Q}^{-\frac{1}{2}}| \\
\text{s.t.} & \quad \text{Trace}(\mathbf{V}_2^{H} \mathbf{V}_2) \leq M_2 \sigma_{\max}.
\end{align*}
\]

(17)

Before solving the optimization problem (OP) in (17), we briefly describe the uniform PA scheme (UPA). The UPA policy can be very useful not only to relax some information assumptions and decrease computational complexity at the transmitter but also because it corresponds to the limit of the optimal PA policy in the high SNR regime.

1) Uniform Power Allocation: In this case, the opportunistic transmitter does not perform any optimization on its own transmit power. It rather uniformly spreads its total power among the previously identified TOs. Thus, the PA matrix \( \mathbf{P}_2 \) is assumed to be of the form

\[
\mathbf{P}_{2,UPA} = \gamma \mathbf{I}_{L_2}.
\]

(18)

where the constant \( \gamma \) is chosen to saturate the transmit power constraint (3),

\[
\gamma = \frac{M_2 \sigma_{\max}}{\text{Trace}(\mathbf{V}_2^{H} \mathbf{V}_2)} = \frac{M_2 \sigma_{\max}}{I_2}.
\]

(19)

2) Optimal Power Allocation: Here, we tackle the OP formulated in (17). For doing so, we assume that the columns of matrix \( \mathbf{V}_2 \) are unitary and mutually orthogonal. We define the matrix \( \mathbf{K} \triangleq \mathbf{Q}^{-(1/2)} \mathbf{H}_2 \mathbf{V}_2 \), where \( \mathbf{K} \) is an \( N_2 \times L_2 \) matrix. Let \( \mathbf{K} = \mathbf{U}_K \mathbf{A}_K \mathbf{V}_K^{H} \) be an SVD of matrix \( \mathbf{K} \), where the matrices \( \mathbf{U}_K \) and \( \mathbf{V}_K \) are unitary matrices with dimensions \( N_2 \times N_2 \) and \( L_2 \times L_2 \) respectively. The matrix \( \mathbf{A}_K \) is an \( N_2 \times L_2 \) matrix with at most \( \min(N_2, L_2) \) nonzero singular values on its main diagonal and zeros in its off-diagonal entries. The entries in the diagonal of the matrix \( \mathbf{A}_K \) are denoted by \( \lambda_{K,1}, \ldots, \lambda_{K, \min(N_2, L_2)} \). Finally, the original OP (17) can be rewritten as

\[
\begin{align*}
\arg \max_{\mathbf{P}_2} & \quad \log_2 |\mathbf{I}_{N_2} + \mathbf{A}_K \mathbf{V}_K^{H} \mathbf{P}_2 \mathbf{V}_K \mathbf{A}_K^{H}| \\
\text{s.t.} & \quad \text{Trace}(\mathbf{P}_2) = \text{Trace}(\mathbf{V}_K^{H} \mathbf{P}_2 \mathbf{V}_K) \leq M_2 \sigma_{\max}.
\end{align*}
\]

(20)

Here, we define the square matrices of dimension \( L_2 \),

\[
\mathbf{P}_2 \triangleq \mathbf{V}_K^{H} \mathbf{P}_2 \mathbf{V}_K,
\]

(21)
and \( \Lambda_K^{u_K} \triangleq \Lambda_K^H \Lambda_K = \text{diag}(\lambda_{K^{u_K}1}, \ldots, \lambda_{K^{u_K}L_2}) \). Using the new variables \( \tilde{P}_2 \) and \( \Lambda_K^{u_K} \), we can write that

\[
\left| I_{N_2} + \Lambda_K V_K^H \tilde{P}_2 V_K \Lambda_K^H \right| = \left| I_{L_2} + \Lambda_K^{u_K} \tilde{P}_2 \right| \\
\leq \prod_{n=1}^{L_2} (1 + \lambda_{K^{u_K}n} \tilde{p}_{2n})
\]

(22)

where \( \tilde{p}_{2n} \), with \( n \in \{1, \ldots, L_2\} \) are the entries of the main diagonal of matrix \( \tilde{P}_2 \). Note that in (22) equality holds if \( \tilde{P}_2 \) is a diagonal matrix [20]. Thus, choosing \( \tilde{P}_2 \) to be diagonal maximizes the transmission rate. Hence, the OP simplifies to

\[
\max_{\tilde{p}_{21}, \ldots, \tilde{p}_{2L_2}} \sum_{n=1}^{L_2} \log_2 (1 + \lambda_{K^{u_K}n} \tilde{p}_{2n})
\]

\[
\text{st.} \quad \sum_{n=1}^{L_2} \tilde{p}_{2n} \leq M_2 \tilde{p}_{2 \max}.
\]

(23)

The simplified optimization problem (23) has eventually a water-filling solution of the form

\[
\forall n \in \{1, \ldots, L_2\}, \quad \tilde{p}_{2n} = \left( \beta_2 - \frac{1}{\lambda_{K^{u_K}n}} \right)^+
\]

(24)

where the water-level \( \beta_2 \) is determined to saturate the power constraints in the optimization problem (23). Once the matrix \( \tilde{P}_2 \) (21) has been obtained using water-filling (24), we define the optimal PA matrix \( \tilde{P}_{2, \text{OPA}} \) by

\[
\tilde{P}_{2, \text{OPA}} = \text{diag}(\tilde{p}_{2,1}, \ldots, \tilde{p}_{2,L_2})
\]

(25)

while the left and right hand factors, \( V_K \) and \( V_K^H \), of matrix \( \tilde{P}_2 \) in (21) are included in the pre-processing matrix:

\[
V_{2, \text{OPA}} = V_2 V_K.
\]

(26)

In the next section, we study the achievable transmission rates of the opportunistic link.

IV. ASYMPTOTIC PERFORMANCE OF THE SECONDARY LINK

In this section, the performance of the secondary link is analyzed in the regime of large number of antennas, which is defined as follows.

**Definition 4 (Regime of Large Numbers of Antennas):** The regime of large numbers of antennas (RLNA) is defined as follows:

- \( \forall i \in \{1, 2\}, N_i \to +\infty \);
- \( \forall j \in \{1, 2\}, M_j \to +\infty \);
- \( \forall (i, j) \in \{1, 2\}^2, \lim_{N_i \to +\infty} (M_j / N_i) = \alpha_{ij} < +\infty \), and \( \alpha_{ij} > 0 \) is constant.

A. Asymptotic Number of Transmit Opportunities

In Section III, two relevant parameters regarding the performance of the opportunistic system can be identified: the number of TOs (\( S \)) and the number of transmit dimensions to which the secondary user can allocate power without affecting the performance of the primary user (\( L_2 \)). Indeed, \( L_2 \) is equivalent to the number of independent symbols the opportunistic system is able to simultaneously transmit. In the following, we analyze both parameters \( S \) and \( L_2 \) in the RLNA by studying the fractions

\[
S_\infty \triangleq \lim_{M_1 \to +\infty} \frac{S}{M_1}
\]

and

\[
L_{2,\infty} \triangleq \lim_{M_1 \to +\infty} \frac{L_2}{M_2}
\]

(27)

(28)

Using (7), the fraction \( S_\infty \) can be re-written as follows

\[
S_\infty = \left( \frac{1}{\alpha_{11}} - m_{1,\infty} \right)^+
\]

(29)

where

\[
m_{1,\infty} \triangleq \lim_{M_1 \to +\infty} \frac{m_1}{M_1}.
\]

(30)

As a preliminary step toward determining the expressions of \( S_\infty \) and \( L_{2,\infty} \), we first show how to find the asymptotic water-level \( \beta_{2,\infty} \) in the RLNA, and the expression of \( m_{1,\infty} \). First, recall from the water-filling solution (4) and the power constraint (3) that

\[
\frac{1}{M_1} \sum_{n=1}^{M_1} p_{1,n} = \frac{1}{M_1} \sum_{n=1}^{M_1} \left( \beta - \frac{\sigma_1^2}{\lambda_{H_{11}^H H_{11}}(n)} \right)^+.
\]

(31)

Define the real function \( q \) by

\[
q(\lambda) = \begin{cases} 0, & \text{if } \lambda = 0, \\ \left( \beta - \frac{\sigma_1^2}{\lambda} \right)^+, & \text{if } \lambda > 0, \end{cases}
\]

(32)

which is continuous and bounded on \( \mathbb{R}^+ \). (31) can be rewritten as

\[
\frac{1}{M_1} \sum_{n=1}^{M_1} q \left( \lambda_{H_{11}^H H_{11}}(n) \right) = \int_{-\infty}^{\infty} q(\lambda) f_{H_{11}^H H_{11}}^M(\lambda) d\lambda
\]

(33)

where \( f_{H_{11}^H H_{11}}^M \) is the probability density function associated with the empirical eigenvalue distribution \( f_{H_{11}^H H_{11}}^M \) of matrix \( H_{11}^H H_{11} \). In the RLNA, the empirical eigenvalue distribution \( f_{H_{11}^H H_{11}}^M \) converges almost surely to the deterministic limiting eigenvalue distribution \( F_{H_{11}^H H_{11}} \), known as the Marčenko–Pastur law [21] whose associated density is

\[
f_{H_{11}^H H_{11}}(\lambda) = \left( 1 - \frac{1}{\alpha_{11}} \right)^+ \delta(\lambda) + \frac{\sqrt{\lambda - a} + \sqrt{b - \lambda}^+}{2\pi \lambda}
\]

(34)

where \( a = (1 - 1/\sqrt{\alpha_{11}})^2 \) and \( b = (1 + 1/\sqrt{\alpha_{11}})^2 \). Note that the Marčenko–Pastur law has a bounded real positive support \( \{0\} \cup [a, b] \) and \( q \) is continuous and bounded on \( \mathbb{R}^+ \).

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Consequently, in the RLNA, we have the almost sure convergence of (33), i.e.,

\[
\int_{-\infty}^{\infty} q(\lambda)f_{HT^{H}_{\hat{H}_{1}}}(\lambda)d\lambda \xrightarrow{a.s.} \int_{-\infty}^{\infty} q(\lambda)f_{HT^{H}_{\hat{H}_{1}}}(\lambda)d\lambda.
\]

Thus, in the RLNA (Def. 4), the water-level \(\beta_{\infty}\) is the unique solution [22] to the equation

\[
\max_{\alpha_{\infty}^{2}} \left( \beta - \frac{\sigma_{\infty}^{2}}{\lambda} \right) \frac{\sqrt{(\lambda - \alpha)(\beta - \lambda)}}{2\pi \lambda} d\lambda - p_{1,\text{max}} = 0,
\]

and it does not depend on any specific realization of the channel transfer matrix \(\hat{H}_{11}\), but only on the maximum power \(p_{1,\text{max}}\) and the receiver noise power \(\sigma_{\infty}^{2}\).

We can now derive \(m_{1,\infty}\). From (5), we have

\[
m_{1,\infty} = \lim_{M_{1} \to \infty} \frac{1}{M_{1}} \sum_{n=1}^{M_{1}} f_{HT_{n}^{H}H_{11}^{n}}(\lambda_{H_{11}^{n}H_{11}^{n}}) = \lim_{M_{1} \to \infty} \frac{1}{M_{1}} \int_{-\infty}^{\infty} \left( \frac{\sigma_{\infty}^{2}}{\lambda} \right) f_{HT_{n}^{H}H_{11}^{n}}(\lambda)d\lambda \xrightarrow{a.s.} \int_{-\infty}^{b} \frac{\sqrt{(\lambda - \alpha)(\beta - \lambda)}}{2\pi \lambda} d\lambda.
\]

Thus, given the asymptotic number of transmit dimensions used by the primary link per primary transmit antenna \(m_{1,\infty}\), we obtain the asymptotic number of transmit opportunities per primary transmit antenna \(S_{\infty}\) by following (27), i.e.,

\[
S_{\infty} = \frac{1}{\alpha_{11}} - \frac{b}{\max_{\alpha_{\infty}^{2}}} \frac{\sqrt{(\lambda - \alpha)(\beta - \lambda)}}{2\pi \lambda} d\lambda.
\]

From (8), the following bounds on \(S_{\infty}\) hold in the RLNA:

\[
\left( \frac{1}{\alpha_{11}} - 1 \right)^{+} \leq S_{\infty} \leq \frac{1}{\alpha_{12}}.
\]

Finally, we give the expression of \(L_{2,\infty}\). Recall that \(L_{2} = \dim \ker(\hat{H}_{1}) = M_{2} - \text{rank}(\hat{H}_{1})\). The rank of \(\hat{H}_{1}\) is given by its number of nonzero singular values, or equivalently by the number of nonzero eigenvalues of matrix \(\hat{H}_{1}^{H}\hat{H}_{1}\). Let \(\lambda_{\hat{H}_{1}^{H}\hat{H}_{1}}^{0}, \ldots, \lambda_{\hat{H}_{1}^{H}\hat{H}_{1}}^{M_{2}}\) denote the eigenvalues of matrix \(\hat{H}_{1}^{H}\hat{H}_{1}\). We have

\[
\frac{1}{\alpha_{11}} - \lim_{N_{1},M_{2} \to \infty} \frac{\text{rank}(\hat{H}_{1})}{M_{2}} = \frac{1}{\alpha_{11}} - \lim_{N_{1},M_{2} \to \infty} \frac{1}{M_{2}} \sum_{n=1}^{M_{2}} \mathbb{1}_{0,\infty}(\lambda_{\hat{H}_{1}^{H}\hat{H}_{1}}^{n}) = L_{2,\infty} \left( \frac{\sigma_{\infty}^{2}}{\alpha_{\infty}^{2}} \right).\]

where \(f_{\hat{H}_{1}^{H}\hat{H}_{1}}(\lambda)\) is the probability density function associated with the empirical eigenvalue distribution \(f_{\hat{H}_{1}^{H}\hat{H}_{1}}(\lambda)\). \(\hat{H}_{1}\) is of size \(M_{1} \times M_{2}\), and the ratio \(M_{2}/M_{1}\) converges in the RLNA to

\[
\tilde{\alpha}_{11} \triangleq \lim_{N_{1},M_{2} \to \infty} \frac{M_{2}}{M_{1}} = \frac{\alpha_{12}}{\alpha_{11}^{2}M_{1,\infty}} < \infty.
\]

Thus, in the RLNA, the empirical eigenvalue distribution \(f_{\hat{H}_{1}^{H}\hat{H}_{1}}(\lambda)\) converges almost surely to the Marčenko-Pastur law [21] \(f_{\hat{H}_{1}^{H}\hat{H}_{1}}(\lambda)\) with associated density

\[
f_{\hat{H}_{1}^{H}\hat{H}_{1}}(\lambda) = \left( \frac{1}{\tilde{\alpha}_{11}^{+}} \right) + \left( \frac{\sqrt{(\lambda - \tilde{c})(\lambda - \tilde{d})}}{2\pi \lambda} \right),
\]

where \(c = \left( \frac{1}{\tilde{\alpha}_{11}^{+}} \right)^{2}\) and \(d = \left( 1 + \frac{1}{\tilde{\alpha}_{11}^{+}} \right)^{2}\). Using (41) in (39) yields

\[
L_{2,\infty} \xrightarrow{a.s.} \frac{1}{\alpha_{11}^{+}} - \frac{b}{\max_{\alpha_{\infty}^{2}}} \frac{\sqrt{(\lambda - \alpha)(\beta - \lambda)}}{2\pi \lambda} d\lambda.
\]

Thus, given the asymptotic water-level \(\beta_{\infty}\) for the primary link, the asymptotic number of TOs per transmit antenna is given by the following expression:

\[
L_{2,\infty} = \left( \frac{1}{\alpha_{11}} - \frac{\alpha_{12}^{+}}{\alpha_{12}^{+} m_{1,\infty}} \right)^{+} = \left( 1 - \frac{1}{\tilde{\alpha}_{11}^{+}} \right)^{+} \left( \frac{1}{\alpha_{11}^{+}} - \frac{\alpha_{12}^{+}}{\alpha_{12}^{+} m_{1,\infty}} \right)^{+}
\]

Note that the number \(S\) of TOs as well as the number \(L_{2}\) of independent symbols that the secondary link can simultaneously transmit are basically determined by the number of antennas and the SNR of the primary system. From (27), it becomes clear that the higher the SNR of the primary link, the lower the number of TOs. Nonetheless, as we shall see in the numerical examples in Section V, for practical values of SNR there exist a nonzero number of TOs the secondary can always exploit.

B. Asymptotic Transmission Rate of the Opportunistic Link

In this subsection, we analyze the behavior of the opportunistic rate per antenna

\[
\tilde{R}_{2}(P_{2},\sigma_{\infty}^{2}) = \frac{1}{N_{2}} \log_{2} \left( N_{2} + Q^{-1} H_{2}^{H} V_{2} P_{2} V_{2}^{H} H_{2}^{H} \right)
\]

(44)
in the RLNA. Interestingly, this quantity can be shown to converge to a limit, the latter being independent of the realization of $H_{22}$. In the present work, we essentially use this limit to conduct a performance analysis of the system under investigation but it is important to know that it can be further exploited, for instance, to prove some properties, or simplify optimization problems [23]. A key transform for analyzing quantities associated with large systems is the Stieltjes transform, which we define in Appendix B. By exploiting the Stieltjes transform and results from random matrix theory for large systems (See Appendix B), it is possible to find the limit of (44) in the RLNA. The corresponding result is as follows.

**Proposition 5 (Asymptotic Transmission Rate):** Define the matrices

$$M_1 \triangleq H_{22} V_{H_{11}} P_1 V_{H_{11}}^P H_{22}^H$$

$$M_2 \triangleq H_{22} V_{2} P_2 V_{2}^P H_{22}^H$$

$$M \triangleq M_1 + M_2,$$

and consider the system model described in Section II with a primary link using the configuration $(V_1, D_1, P_1)$ described in Section III-A, and a secondary link with the configuration $(V_2, D_2, P_2)$ described in Sections III-C, III-D, with $P_2$ any PA matrix independent from the noise level $\sigma_2^2$. Then, in the RLNA (Def. 4), under the assumption that $P_1$ and $V_2 P_2 V_2^H$ have limiting eigenvalue distributions $F_{P_1}$ and $F_{V_2 P_2 V_2^H}$ with compact support, the transmission rate per antenna of the opportunistic link $(T_{X_2} - R_{X_2})$ converges almost surely to

$$R_{2,\infty} = \frac{1}{\ln(2)} \int_0^{+\infty} \frac{G_M(-z) - G_M(-z)}{\sigma_2^2} dz$$

(48)

where $G_M(z)$ and $G_M(z)$ are the Stieltjes transforms of the limiting eigenvalue distribution of matrices $M_1$ and $M_2$, respectively. $G_M(z)$ and $G_M(z)$ are obtained by solving the fixed point equations (with unique solution when $z \in \mathbb{R}^+$) [24]:

$$G_{M_1}(z) = \frac{-1}{z - g(G_{M_1}(z))}$$

(49)

and

$$G_{M_2}(z) = \frac{-1}{z - g(G_{M_2}(z)) - h(G_{M_2}(z))}$$

(50)

respectively, where the functions $g(u)$ and $h(u)$ are defined as follows:

$$g(u) \triangleq \mathbb{E}

\left[

\frac{p_1}{1 + \frac{1}{\alpha_2^2} p_1 u}

\right]$$

(51)

$$h(u) \triangleq \mathbb{E}

\left[

\frac{p_2}{1 + \frac{1}{\alpha_2^2} p_2 u}

\right]$$

(52)

with the expectations in (51) and (52) taken on the random variables $p_1$ and $p_2$ with distribution $F_{P_1}$ and $F_{V_2 P_2 V_2^H}$, respectively.

**Proof:** For the proof, see Appendix C.

The (non trivial) result in Proposition 5 holds for any power allocation matrix $P_2$ independent of $\sigma_2^2$. In particular, the case of the uniform power allocation policy perfectly meets this assumption. This also means that it holds for the optimum PA policy in the high SNR regime. For low and medium SNRs, the authors have noticed that the matrix $P_{2,\text{OPA}}$ is in general not independent of $\sigma_2^2$. This is because $P_2$ is obtained from a water-filling procedure. The corresponding technical problem is not trivial and is therefore left as an extension of the present work.

V. NUMERICAL RESULTS

A. The Number $S$ of Transmit Opportunities

As shown in (27), the number of TOs is a function of the number of antennas and the SNR of the primary link. In Fig. 2, we plot the number of TOs per transmit antenna $S_{X_2}$ as a function of the SNR for different number of antennas in the receiver and transmitter of the primary link. Interestingly, even though the number of TOs is a nonincreasing function of the SNR, Fig. 2 shows that for practical values of the SNR (10-20 dBs) there exists a nonzero number of TOs. Note also that the number of TOs is an increasing function of the ratio $(\alpha_{11} \sim \alpha_{11} / N_1)$. For instance, in the case $N_1 > M_1$, i.e., $\alpha_{11} > 1$ the secondary transmitters always see a nonzero number of TOs independently of the SNR of the primary link, and thus, opportunistic communications are always feasible. On the contrary, when $\alpha_{11} \leq 1$, the feasibility of opportunistic communications depends on the SNR of the primary link.

Finally, it is important to remark that even though, the analysis of the number of TOs has been done in the RLNA (Def. 4), the model is also valid for finite number of antennas. In Fig. 2, we have also plotted the number of TOs observed for a given realization of the channel transfer matrix $H_{11}$ when $N_1 = 10$ and $\alpha_{11} \in \{(1/2), 1, 2\}$. Therein, it can be seen how the theoretical result from (27) matches the simulation results.
B. Comparison Between OIA and ZFBF

We compare our OIA scheme with the zero-forcing beamforming (ZFBF) scheme [18]. Within this scheme, the pre-processing matrix $\mathbf{V}_2$, denoted by $\mathbf{V}_{2,ZFBF}$, satisfies the condition

$$\mathbf{H}_{12} \mathbf{V}_{2,ZFBF} = \mathbf{0}_{N_r,N_L_2}$$

(53)

which implies that ZFBF is feasible only in some particular cases regarding the rank of matrix $\mathbf{H}_{12}$. For instance, when $M_2 \leq N_1$ and $\mathbf{H}_{12}$ is full column rank, the pre-processing matrix is the null matrix, i.e., $\mathbf{V}_{2,ZFBF} = \mathbf{0}_{M_2,N_2}$ and thus, no transmission takes place. On the contrary, in the case of OIA when $M_2 \leq N_1$, it is still possible to opportunistically transmit with a nonnull matrix $\mathbf{V}_2$ in two cases as shown in Section III-C:

- if $m_1 < M_2$;
- or if $m_1 \geq M_2$ and $\mathbf{H}_1$ is not full column rank.

Another remark is that when using ZFBF and both primary and secondary receivers come close, the opportunistic link will observe a significant power reduction since both the targeted and nulling directions become difficult to distinguish. This power reduction will be less significant in the case of OIA since it always holds that $\text{rank}(\mathbf{V}_2) \geq \text{rank}(\mathbf{V}_{2,ZFBF})$ thanks to the existence of the additional TOs. Strict equality holds only when $S = ((1/\alpha_{\text{OPA}}) - 1)$. As discussed in Section III-B, the number of TOs ($S$) is independent of the position of one receiver with respect to the other. It rather depends on the channel realization $\mathbf{H}_{11}$ and the SNR of the primary link.

In the following, for the ease of presentation, we consider that both primary and secondary devices are equipped with the same number of antennas $N_s = N_1 = N_2$ and $N_r = M_1 = M_2$, respectively. In this scenario, we consider the cases where $N_r > N_r$ and $N_r \leq N_r$.

1) Case $N_r > N_r$: In Fig. 3, we consider the case where $\alpha \approx 5/4$, with $N_r \in \{3, 9\}$. In this case, we observe that even for a small number of antennas, the OIA technique is superior to the classical ZFBF. Moreover, the higher the number of antennas, the higher the difference between the performance of both techniques. An important remark here is that, at high SNR, the performance of ZFBF and OIA is almost identical. This is basically because at high SNR, the number of TOs tends to its lower bound $N_r - N_r$ (from (8)), which coincides with the number of spatial directions to which ZFBF can avoid interfering. Another remark is that both UPA and OPA schemes perform identically at high SNR.

2) Case $N_r \leq N_r$: In this case, the ZFBF solution is not feasible and thus, we focus only on the OIA solution. In Fig. 4, we plot the transmission rate for the case where $N_r = N_r \in \{3, 6, 9\}$. We observe that at high SNR for the primary link and small number of antennas, the uniform PA performs similarly as the optimal PA. For a higher number of antennas and low SNR in the primary link, the difference between the uniform and optimal PA is significant. To show the impact of the SINR of both primary and secondary links on the opportunistic transmission rate, we present Fig. 5. Therein, it can be seen clearly that the transmission rate in the opportunistic link is inversely proportional to the SINR level at the primary link. This is due to the lack of TOs as stated in Section III-B. For the case when $N_r < N_r$ with strict inequality, an opportunistic transmission takes place only if $N_r - N_r \leq S$ and $\mathbf{H}_{11}$ is not full column rank. Here, the behavior of the opportunistic transmission rate is similar to the case $N_r = N_r$ with the particularity that the opportunistic transmission rate reaches zero at a lower SNR level. As in the previous case, this is also a consequence of the number of available TOs.

C. Asymptotic Transmission Rate

In Fig. 6, we plot both primary and secondary transmission rates for a given realization of matrices $\mathbf{H}_{(i,j)}$ $\forall (i,j) \in \{1,2\}$

(52)

We also plot the asymptotes obtained from Proposition 5 considering UPA in the secondary link and the optimal PA of the primary link (4). We observe that in both cases the transmission rate converges rapidly to the asymptotes even for a small number of antennas. This shows that Proposition 5 constitutes a good estimation of the achievable transmission rate for the secondary link even for finite number of antennas. We use Proposition 5 to compare the asymptotic transmission rate of the secondary and primary link. The asymptotic transmission rate of the primary receiver corresponds to the capacity of a single user $N_r \times N_r$ MIMO link whose asymptotes are provided in [25]. From Fig. 6, it becomes evident how the secondary link is able
to achieve transmission rates of the same order as the primary link depending on both its own SNR and that of the primary link.

VI. CONCLUSION

In this paper, we proposed a technique to recycle spatial directions left unused by a primary MIMO link, so that they can be reused by secondary links. Interestingly, the number of spatial directions can be evaluated analytically and shown to be sufficiently high to allow a secondary system to achieve a significant transmission rate. We provided a signal construction technique to exploit those spatial resources and a power allocation policy which maximizes the opportunistic transmission rate. Based on our asymptotical analysis, we show that this technique allows a secondary link to achieve transmission rates of the same order as those of the primary link, depending on their respective SNRs. To mention few interesting extensions of this work, we recall that our solution concerns only two MIMO links. The case where there exists several opportunistic devices and/or several primary devices remains to be studied in details. More importantly, some information assumptions could be relaxed to make the proposed approach more practical. This remark concerns CSI assumptions but also behavioral assumptions. Indeed, it was assumed that the precoding scheme used by the primary transmitter is capacity-achieving, which allows the secondary transmitter to predict how the secondary transmitter is going to exploit its spatial resources. This behavioral assumption could be relaxed but some spatial sensing mechanisms should be designed to know which spatial modes can be effectively used by the secondary transmitter, which could be an interesting extension of the proposed scheme.

APPENDIX A

PROOF OF LEMMA 1

Here, we prove Lemma 1 which states that: if a matrix $\mathbf{V}_2$ satisfies the condition $\mathbf{H}_1\mathbf{V}_2 = 0_{(N_1 - S) \times L_2}$ then it meets the IA condition (3).

Proof: Let $\mathbf{H}_{11} = \mathbf{U}_{H11}\mathbf{A}_{H11}\mathbf{V}_{H11}^H$ be a sorted SVD of matrix $\mathbf{H}_{11}$, with $\mathbf{U}_{H11}$ and $\mathbf{V}_{H11}$ two unitary matrices of sizes $N_1 \times N_1$ and $M_1 \times M_1$, respectively, and $\mathbf{A}_{H11}$ an $N_1 \times M_1$ matrix with main diagonal $(\lambda_{H11,1}, \ldots, \lambda_{H11,\min(N_1,M_1)})$ and zeros on its off-diagonal, such that $\lambda_{H11,1} \geq \lambda_{H11,2} \geq \cdots \geq \lambda_{H11,\min(N_1,M_1)}$. Given that the singular values of the matrix $\mathbf{H}_{11}$ are sorted, we can write matrix $\mathbf{A}_{H11}\mathbf{P}_1\mathbf{A}_{H11}^H$ as a block matrix,

$$\mathbf{A}_{H11}\mathbf{P}_1\mathbf{A}_{H11}^H = \begin{pmatrix} \mathbf{0}_{(N_1 - m_1)\times m_1} & \mathbf{0}_{(N_1 - m_1)\times (N_1 - m_1)} \\ \mathbf{0}_{(N_1 - m_1)\times m_1} & \mathbf{0}_{(N_1 - m_1)\times (N_1 - m_1)} \end{pmatrix}$$

where the diagonal matrix $\mathbf{\Psi}$ of size $m_1 \times m_1$ is $\mathbf{\Psi} = \text{diag}(\lambda_{H11,1}^2, \lambda_{H11,2}^2, \ldots, \lambda_{H11,m_1}^2)$.

Now let us split the interference-plus-noise covariance matrix (9) as

$$\mathbf{R} = \begin{pmatrix} m_1 & N_1 - m_1 \\ N_1 - m_1 & m_1 \end{pmatrix} \begin{pmatrix} \mathbf{R}_{1} + \sigma_1^2 \mathbf{I}_{m_1} & \mathbf{R}_{2} \\ \mathbf{R}_{2}^H & \mathbf{R}_{3} + \sigma_3^2 \mathbf{I}_{N_1 - m_1} \end{pmatrix},$$

where $(\mathbf{R}_{4} + \sigma_1^2 \mathbf{I}_{m_1})$ and $(\mathbf{R}_{3} + \sigma_3^2 \mathbf{I}_{N_1 - m_1})$ are invertible Hermitian matrices, and matrices $\mathbf{R}_{4}, \mathbf{R}_{2}$ and $\mathbf{R}_{3}$ are defined from (9) and (11) as

$$\mathbf{R}_{1} = \mathbf{H}_1\mathbf{V}_2\mathbf{P}_2\mathbf{V}_2^H\mathbf{H}_1^H$$

$$\mathbf{R}_{2} = \mathbf{H}_1\mathbf{V}_2\mathbf{P}_2\mathbf{V}_2^H\mathbf{H}_1^H$$

$$\mathbf{R}_{3} = \mathbf{H}_1\mathbf{V}_2\mathbf{P}_2\mathbf{V}_2^H.$$

Now, by plugging expressions (54) and (55) in (10), the IA condition can be rewritten as follows:

$$\log_2 \left| \sigma_1^2 \mathbf{I}_{m_1} + \mathbf{\Psi} \right| - \log_2 \left| \sigma_1^2 \mathbf{I}_{N_1} \right| = \log_2 \left| \mathbf{R}_{1} + \sigma_1^2 \mathbf{I}_{m_1} + \mathbf{\Psi} \right| - \log_2 \left| \mathbf{R}_{1} + \sigma_1^2 \mathbf{I}_{m_1} \right|$$

$$- \log_2 \left| \mathbf{R}_{3} + \sigma_3^2 \mathbf{I}_{N_1 - m_1} - \mathbf{R}_{2}^H \mathbf{R}_{1} + \sigma_1^2 \mathbf{I}_{m_1} + \mathbf{\Psi} \right| - \mathbf{R}_{1} \mathbf{R}_{2}$$

$$= \log_2 \left| \mathbf{R}_{4} + \sigma_1^2 \mathbf{I}_{m_1} \right| - \log_2 \left| \mathbf{R}_{3} + \sigma_3^2 \mathbf{I}_{N_1 - m_1} - \mathbf{R}_{2}^H (\mathbf{R}_{1} + \sigma_1^2 \mathbf{I}_{m_1} + \mathbf{\Psi})^{-1} \mathbf{R}_{2} \right|.$$

Note that there exists several choices for the submatrices $\mathbf{R}_{1}, \mathbf{R}_{2},$ and $\mathbf{R}_{3}$ allowing the equality in (59) to be met. We see that
a possible choice in order to meet the IA condition is $R_1 = 0$, $R_2 = 0$, independently of the matrix $R_3$. Thus, from (56) and (57) we have $R_3 = 0$ and $R_2 = 0$ by imposing the condition $H_1 V_2 = 0_{m_1 \times L_2}$, for any given PA matrix $P_2$, which concludes the proof.

**APPENDIX B**

**DEFINITIONS**

In this appendix, we present useful definitions and previous results used in the proofs of Appendix C.

**Definition 6:** Let $X$ be an $n \times n$ random matrix with empirical eigenvalue distribution function $F_{X}^{(n)}$. We define the following transforms associated with the distribution $F_{X}^{(n)}$, for $z \in \mathbb{C}^+$ such that $\text{Im}(z) > 0$:

Stieltjes transform $G_X(z) \triangleq \int_{-\infty}^{\infty} \frac{1}{t-z} dF_{X}^{(n)}(t)$ \hspace{1cm} (60)

$\Upsilon_X(z) \triangleq \int_{-\infty}^{\infty} \frac{zt}{1-zt} dF_{X}^{(n)}(t)$ \hspace{1cm} (61)

S-transform $S_X(z) \triangleq \int_{-\infty}^{\infty} \frac{1+z}{z} \Upsilon_X^{-1}(z)$ \hspace{1cm} (62)

where the function $\Upsilon_X^{-1}(z)$ is the reciprocal function of $\Upsilon_X(z)$, i.e.,

$\Upsilon_X^{-1}(\Upsilon_X(z)) = \Upsilon_X(\Upsilon_X^{-1}(z)) = z$. \hspace{1cm} (63)

From (60) and (61), we obtain the following relationship between the function $\Upsilon_X(z)$ (named $\Upsilon$-transform in [26]) and the Stieltjes transform $G_X(z)$.

$\Upsilon_X(z) = -1 - \frac{1}{z} G_X \left( \frac{1}{z} \right).$ \hspace{1cm} (64)

**APPENDIX C**

**PROOF OF PROPOSITION 5**

In this appendix, we provide a proof of Proposition 5 on the asymptotic expression of the opportunistic transmission rate per antenna, defined by

$R_{2,\infty}(P_2, \sigma_2^2) \triangleq \lim_{\gamma \to \infty} \lim_{N_2, M \to \infty} R_2(P_2, \sigma_2^2)$, \hspace{1cm} (65)

First, we list the steps of the proof and then we present a detailed development for each of them:

**Step 1:** Express $\frac{\partial R_{2,\infty}(P_2, \sigma_2^2)}{\partial \sigma_2^2}$ as a function of the Stieltjes transforms $G_{M_1}(z)$ and $G_{M}(z)$.

Step 2: Obtain $G_{M_1}(z)$.

Step 3: Obtain $G_{M}(z)$.

Step 4: Integrate $\frac{\partial R_{2,\infty}(P_2, \sigma_2^2)}{\partial \sigma_2^2}$ to obtain $R_{2,\infty}(P_2, \sigma_2^2)$.

**Step 1:** Express $\frac{\partial R_{2,\infty}(P_2, \sigma_2^2)}{\partial \sigma_2^2}$ as a function of the Stieltjes Transforms $G_{M_1}(z)$ and $G_{M}(z)$: Using (16) and (15), the opportunistic rate per receive antenna $R_2$ can be rewritten as follows:

$R_2(P_2, \sigma_2^2) = \frac{1}{N_2} \log_2 \left| I_{N_2} + Q^{-\frac{1}{2}} H_{22} V_2 P_2 V_2^H H_{22}^H Q^{-\frac{1}{2}} \right|$

$= \frac{1}{N_2} \log_2 \left| \sigma_2^2 I_{N_2} + M_1 + M_2 \right|$

$- \frac{1}{N_2} \log_2 \left| \sigma_2^2 I_{N_2} + M_1 \right|$ \hspace{1cm} (66)

with $M_1 \triangleq H_{22} V_2 P_2 V_2^H H_{22}^H$, and $M = M_1 + M_2$. Matrices $M$ and $M_1$ are Hermitian Gramian matrices with eigenvalue decomposition $M = U_M \Lambda_M U_M^H$ and $M_1 = U_M \Lambda_M I_M^H$, respectively. Matrix $U_M$ and $U_M$ are $N_2 \times N_2$ unitary matrices, and $\Lambda_M = \text{diag} (\lambda_{M,1}, \ldots, \lambda_{M,N_2})$ and $\Lambda_M = \text{diag} (\lambda_{1,1}, \ldots, \lambda_{N_2,N_2})$ are square diagonal matrices containing the eigenvalues of the matrices $M$ and $M_1$ in decreasing order. Expression (65) can be written as

$R_2(P_2, \sigma_2^2) = \frac{1}{N_2} \sum_{\lambda \in \mathbb{C}} \log_2 \left( \sigma_2^2 + \lambda M_2 \right)$

$- \log_2 \left( \sigma_2^2 + \lambda M_1 \right)$

$= \int \log_2 (\lambda + \sigma_2^2) dF_{N_2}(\lambda)$

$- \log_2 (\lambda + \sigma_2^2) dF_{M}(\lambda)$

$= \int \log_2 (\lambda + \sigma_2^2) dF_{M}(\lambda)$

$- \int \log_2 (\lambda + \sigma_2^2) dF_{M}(\lambda)$ \hspace{1cm} (66)

where $F_{N_2}$ and $F_{M_1}$ are respectively the empirical eigenvalue distributions of matrices $M$ and $M_1$ of size $N_2$, that converge almost surely to the asymptotic eigenvalue distributions $F_M$ and $F_{M_1}$, respectively. $F_M$ and $F_{M_1}$ have a compact support. Indeed the empirical eigenvalue distribution of Wishart matrices $H_{ij} H_{ij}^H$ converges almost surely to the compactly supported Marcenko-Pastur law, and by assumption, matrices $V, P, V_i^H, i \in \{1, 2\}$ have a limit eigenvalue distribution with a compact support. Then by Lemma 5 in [27], the asymptotic eigenvalue distribution of $M_1$ and $M_2$ have a compact support. The logarithm function being continuous, it is bounded on the compact supports of the asymptotic eigenvalue distributions of $M_1$ and $M$, therefore, the almost sure convergence in (66) could be obtained by using the bounded convergence theorem [28].

From (66), the derivative of the asymptotic rate $R_{2,\infty}(P_2, \sigma_2^2)$ with respect to the noise power $\sigma_2^2$ can be written as

$\frac{\partial}{\partial \sigma_2^2} R_{2,\infty}(P_2, \sigma_2^2)$

$= \frac{1}{\ln 2 \upsilon} \left( \int \frac{1}{\sigma_2^2 + \lambda} dF_M(\lambda) - \int \frac{1}{\sigma_2^2 + \lambda} dF_{M_1}(\lambda) \right)$

$= \frac{1}{\ln 2 \upsilon} \left( G_M(\sigma_2^2) - G_{M_1}(\sigma_2^2) \right)$ \hspace{1cm} (67)
where $G_M(z)$ and $G_M(z)$ are the Stieltjes transforms of the asymptotic eigenvalue distributions $F_M$ and $F_M$, respectively.

**Step 2: Obtain $G_M(z)$:** Matrix $M_1$ can be written as

$$M_1 = \sqrt{\alpha_{21}} \mathbf{H}_2 \mathbf{V}_{H_1} \frac{\mathbf{P}_1 \mathbf{V}_{H_1}^H \mathbf{H}_{21}^H}{\alpha_{21}}.$$

(68)

The entries of the $N_2 \times M_1$ matrix $\sqrt{\alpha_{21}} \mathbf{H}_2 \mathbf{V}_{H_1}$ are zero-mean i.i.d. complex Gaussian with variance $(\alpha_{21}/M_1) = (1/N_2)$, thus $\sqrt{\alpha_{21}} \mathbf{H}_2 \mathbf{V}_{H_1}$ is bi-unitarily invariant. Matrix $\mathbf{V}_{H_1}$ is unitary, consequently $\sqrt{\alpha_{21}} \mathbf{H}_2 \mathbf{V}_{H_1}$ has the same distribution as $\sqrt{\alpha_{21}} \mathbf{H}_2$, in particular its entries are i.i.d. with mean zero and variance $1/N_2$. From (4), $\mathbf{P}_1/\alpha_{21}$ is diagonal, and by assumption it has a limit eigenvalue distribution $F_{P_1/\alpha_{21}}$. Thus, we can apply Theorem 1.1 in [24] to $M_1$, in the particular case where $\mathbf{A} = \mathbf{0}_{N_2}$ to obtain the Stieltjes transform of the asymptotic eigenvalue distribution of matrix $M_1$:

$$G_M(z) = G_{N_2} \left( z - \frac{\lambda}{1 + \alpha_{21}} \int_0^\infty \frac{dF_{G_{N_2}}(\lambda)}{\lambda} \right)$$

$$= G_{N_2} \left( z - \frac{1}{1 + \alpha_{21}} \int_{-\infty}^\infty \frac{\lambda}{\lambda} dF_{G_{N_2}}(\lambda) \right)$$

$$= G_{N_2} \left( z - g(G_{N_2}(z)) \right).$$

(69)

where the function $g(u)$ is defined by

$$g(u) = \frac{t}{1 + \alpha_{21} u} \int_{-\infty}^{\infty} \frac{t}{1 + \alpha_{21} u} f_{P_1}(t) dt = \mathbb{E}\left[ \frac{t}{1 + \alpha_{21} u} \right],$$

where the random variable $t$ follows the c.d.f. $f_{P_1}$.

The square null matrix $\mathbf{A}$ has an asymptotic eigenvalue distribution $F_\mathbf{0}(\lambda) = \mu(\lambda)$. Thus, its Stieltjes transform is

$$G_\mathbf{A}(\lambda) = \int_{-\infty}^\infty \frac{1}{\lambda - z} d\lambda = -\frac{1}{z}.$$

(70)

Then, using expressions (69) and (70), we obtain

$$G_M(z) = \frac{1}{z - g(G_M(z))}.$$  

(71)

Expression (71) is a fixed-point equation with unique solution when $z \in \mathbb{R}^-$ [24].

**Step 3: Obtain $G_M(z)$:** Recall that

$$M_1 \triangleq H_1 H_2^T V_2 P_2 V_2^H H_2^H + H_1 H_1^T P_1 V_1^H H_1^H.$$

(72)

To obtain the Stieltjes transform $G_M(z)$, we apply Theorem 1.1 in [24] as in Step 2:

$$G_M(z) = G_M(z - g(G_M(z))),$$

(73)

To obtain the Stieltjes transform $G_M(z)$ of the asymptotic eigenvalue distribution function of the matrix $M_2 = H_2 V_2 P_2 V_2^H H_2^H$, we first express its $S$-transform as

$$S_{M_2}(z) = S_{H_2 V_2 P_2 V_2^H H_2^H}(z)$$

$$= \frac{\lambda}{\alpha_{21} \mathbf{H}_2 \mathbf{V}_{H_1} \frac{\mathbf{P}_1 \mathbf{V}_{H_1}^H \mathbf{H}_{21}^H}{\alpha_{21}}}.$$

and by Lemma 1 in [27]:

$$S_{M_2}(z) = \left( \frac{z + 1}{z + 2} \right) \left( \frac{\lambda}{\alpha_{21} \mathbf{H}_2 \mathbf{V}_{H_1} \frac{\mathbf{P}_1 \mathbf{V}_{H_1}^H \mathbf{H}_{21}^H}{\alpha_{21}}} \right)^2,$$

(74)

and by Theorem 1 in [29]:

$$S_{M_2}(z) = \left( \frac{z + 1}{z + \alpha_{21}} \right) \left( \frac{\lambda}{\alpha_{21}} \right)^2$$

$$\times \left( \frac{\lambda}{\alpha_{21}} \right) \left( \frac{\lambda}{\alpha_{21}} \right)^2$$

$$\left( \frac{\lambda}{\alpha_{21}} \right) \left( \frac{\lambda}{\alpha_{21}} \right)^2.$$$$

The $S$-transforms $S_{M_2}(z)$ and $S_{V_2 P_2 V_2^H}(z/\alpha)$ in expression (74) can be written as functions of their $\Upsilon$-transforms:

$$\Upsilon_{M_2}^{-1}(z) = \left( \frac{1}{z + 1} \right) \Upsilon_{V_2 P_2 V_2^H}(\lambda / \alpha_{21})$$

(75)

$$\Upsilon_{V_2 P_2 V_2^H}(\lambda / \alpha_{21}) = \left( \frac{1}{z + 1} \right) \Upsilon_{V_2 P_2 V_2^H}(\lambda / \alpha_{21}).$$

(76)

Then, plugging (75) and (76) into (74) yields

$$\Upsilon_{M_2}^{-1}(z) = \left( \frac{1}{z + 1} \right) \Upsilon_{V_2 P_2 V_2^H}(\lambda / \alpha_{21}).$$

(77)

Now, using the relation (64) between the $\Upsilon$-transform and the Stieltjes transform, we write

$$G_M(z) = \left( \frac{1}{z + 1} \right) \left( \frac{1}{z + g(G_M(z))} \right).$$

(78)

and from (73), we obtain

$$G_M(z) = \left( \frac{1}{z + g(G_M(z))} \right) \left( \frac{1}{z + g(G_M(z))} + 1 \right).$$

(79)

We handle (79) to obtain $G_M(z)$ as a function of $\Upsilon_{V_2 P_2 V_2^H}(z)$:

$$\Upsilon_{M_2}^{-1}(z) = \left( \frac{1}{z + g(G_M(z))} \right) \left( \frac{1}{z + g(G_M(z))} \right) \left( \frac{1}{z + g(G_M(z))} + 1 \right).$$
\[
\frac{1}{z - g \left( G_M(z) \right)} = \frac{-1}{(z - g \left( G_M(z) \right)) G_M(z)} \\
\times \gamma^{\alpha_{22}} \frac{\eta_1}{\alpha_{22} V_2^{\alpha_{22}/2} V_2'} \left( \frac{1 + (z - g \left( G_M(z) \right)) G_M(z)}{\alpha_{22}} \right) \\
- G_M(z) \\
= \gamma^{\alpha_{22}} \frac{\eta_1}{\alpha_{22} V_2^{\alpha_{22}/2} V_2'} \left( \frac{1 + (z - g \left( G_M(z) \right)) G_M(z)}{\alpha_{22}} \right) \gamma^{-1} \frac{\eta_1}{\alpha_{22} V_2^{\alpha_{22}/2} V_2'} (-G_M(z)) \\
= \frac{1}{z - g \left( G_M(z) \right)} \left( 1 + \frac{1}{\alpha_{22} \gamma V_2^{\alpha_{22}/2} V_2'} (-G_M(z)) \right). \\
\tag{80}
\]

From the definition of the $\gamma$-transform (61), it follows that

\[
\alpha_{22} \gamma V_2^{\alpha_{22}/2} V_2' (-G_M(z)) = \alpha_{22} \int \frac{G_M(z) \lambda}{1 + G_M(z) \lambda} dF_{V_2^{\alpha_{22}/2} V_2'}(\lambda) \\
= \int \frac{\alpha_{22} G_M(z) \lambda}{1 + G_M(z) \lambda} f_{V_2^{\alpha_{22}/2} V_2'}(\alpha_{22} \lambda) \lambda d\lambda \\
= \int \frac{-G_M(z) \lambda}{1 + G_M(z) \lambda} f_{V_2^{\alpha_{22}/2} V_2'}(\alpha_{22} \lambda) \lambda dt. \\
\tag{81}
\]

Using (81) in (80), we have

\[
G_M(z) = \frac{1}{z - g \left( G_M(z) \right)} \left( 1 - G_M(z) h \left( G_M(z) \right) \right) \\
\tag{82}
\]

with the function $h(u)$ defined as follows

\[
h(u) = \int \frac{t}{1 + \frac{u}{\alpha_{22}} t} dF_{V_2^{\alpha_{22}/2} V_2'}(t) = E \left[ \frac{p_2}{1 + \frac{p_2}{\alpha_{22}} p_2 u} \right] \\
\]

where the random variable $p_2$ follows the distribution $F_{V_2^{\alpha_{22}/2} V_2'}$. Factorizing $G_M(z)$ in (82) finally yields

\[
G_M(z) = \frac{-1}{z - g \left( G_M(z) \right) - h \left( G_M(z) \right)}. \\
\tag{83}
\]

Expression (83) is a fixed point equation with unique solution when $z \in \mathbb{R}^- \setminus \{0\}$. 

Step 4: Integrate $\partial \tilde{R}_{2,\infty}(\mathbf{p}_2, \sigma_2^2)/\partial \sigma_2^2$ to Obtain $\tilde{R}_{2,\infty}(\mathbf{p}_2, \sigma_2^2)$ in the RLNA: From (67), we have that

\[
\frac{\partial}{\partial \sigma_2^2} \tilde{R}_{2,\infty}(\mathbf{p}_2, \sigma_2^2) = \frac{1}{\ln 2} \left( G_M \left( -\sigma_2^2 \right) - G_M \left( -\sigma_2^2 \right) \right). \\
\tag{84}
\]

Moreover, it is known that if $\sigma_2^2 \to \infty$ no reliable communication is possible and thus, $\tilde{R}_{2,\infty} = 0$. Hence, the asymptotic rate of the opportunistic link can be obtained by integrating expression (84)

\[
\tilde{R}_{\infty} = \frac{-1}{\ln 2} \int_{\sigma_2^2}^{\infty} \left( G_M \left( -\sigma_2^2 \right) - G_M \left( -\sigma_2^2 \right) \right) d\sigma_2^2 \\
\tag{85}
\]

which ends the proof.

REFERENCES


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